

↳ The degree sequence of the graph (e.g.)

■ **Example 4:** Prove that there does not exist a polyhedron with an odd number of faces such that each face has an odd number of edges.

✓ Such a polyhedron does not exist.

Prove:

- ① Suppose there are n faces, and each face has a_1, a_2, \dots, a_n edges.
- ② If all these numbers are odd, then the sum $a_1 + a_2 + \dots + a_n$ is an odd number.
- ③ However, each edge is shared by exactly two faces, which means each edge is counted twice in the sum above.
- ④ Therefore, if there are m edges in total, then $2m = a_1 + a_2 + \dots + a_n$, this leads to an even number equaling an odd number — a **contradiction**.

↳ The degree sequence of the graph (e.g.)

■ **Example 5:** Suppose a simple undirected graph of order 9 has each vertex of degree 5 or 6. Prove that it has **at least 5 vertices of degree 6** or **at least 6 vertices of degree 5**.

Proof 1: Consider all possible cases. Suppose there are a vertices of degree 5 and b vertices of degree 6.

(1) $a=0, b=9$;

(2) $a=2, b=7$;

(3) $a=4, b=5$;

(4) $a=6, b=3$;

(5) $a=8, b=1$

(1)–(3): At least 5 vertices of degree 6

(4) and (5): At least 6 vertices of degree 5

Proof 2: Suppose $b < 5$, then $a > 9 - 5 = 4$.
By a corollary of the Handshaking Lemma, $a \geq 6$.

↳ 6.1 Basic Concepts of Graphs

- 6.1.1 Undirected and Directed Graphs
- 6.1.2 Vertex Degree and the Handshaking Lemma
- 6.1.3 Common Types of Graphs
- 6.1.4 Subgraphs and Complements
- 6.1.5 Graph Isomorphism

↳ Multigraphs and Simple Graphs

■ Definition 6.3:

- (1) In an *undirected graph*, two or more edges connecting the same pair of vertices are called *parallel edges*, and the number of such edges is called the **multiplicity**.
- (2) In a *directed graph*, two or more edges with the same starting and ending vertices are called *directed parallel edges* (or simply **parallel edges**), and their number is also referred to as the **multiplicity**.
- (3) A graph that contains parallel edges is called a *multigraph*.
- (4) A graph that has neither parallel edges nor loops is called a *simple graph*.

↳ Determining a Simple Graph (e.g.)

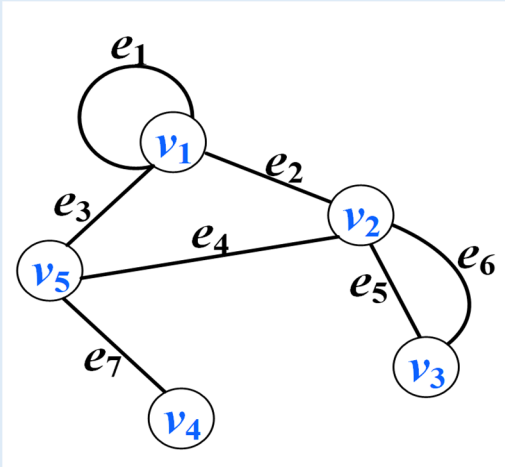


Figure 1

- Edges e_5 and e_6 are parallel edges.
- The multiplicity is 2.
- It is **not a simple graph**.

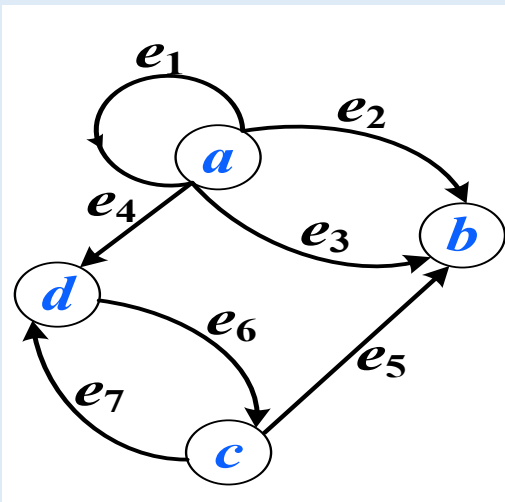
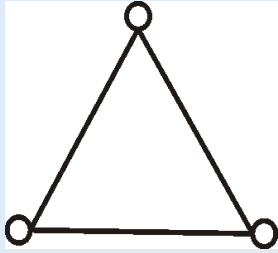
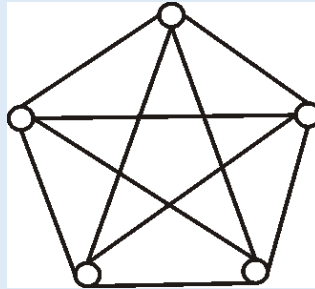
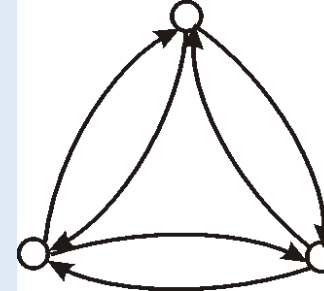
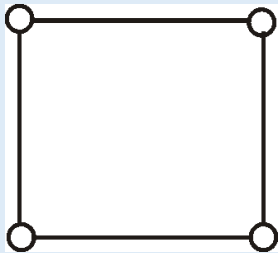
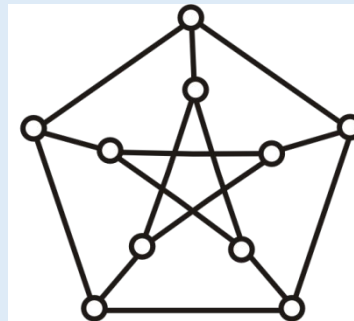
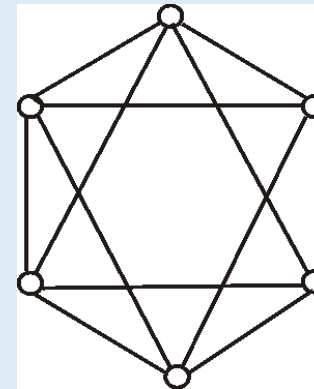


Figure 2

- Edges e_2 and e_3 are parallel edges, with a multiplicity of 2.
- Edges e_6 and e_7 are not parallel edges.
- It is **not a simple graph**.

↳ Complete Graphs and Regular Graphs

- **Undirected complete graph**: An undirected simple graph in which every pair of vertices is connected by an edge.
 - An undirected complete graph of order n is denoted by K_n , Number of vertices: n , Number of edge: $m=n(n-1)/2$, $\Delta=\delta=n-1$.
- **Directed complete graph**: A directed simple graph in which every pair of vertices is connected by two edges in opposite directions.
 - Number of vertices: n , Number of edges: $m=n(n-1)$, $\Delta^+=\delta^+=\Delta^-=\delta^-=n-1$, $\Delta=\delta=2(n-1)$.
- **k -regular graph**: An undirected simple graph in which every vertex has degree k .
 - Number of vertices: n , Number of edges: $m=kn/2$.

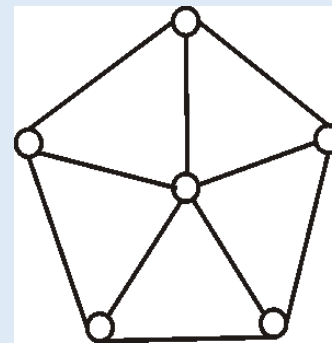
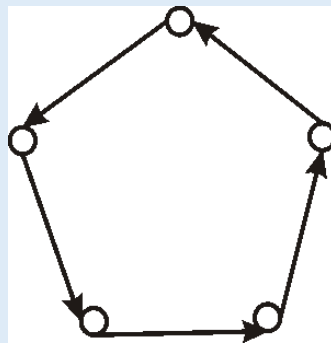
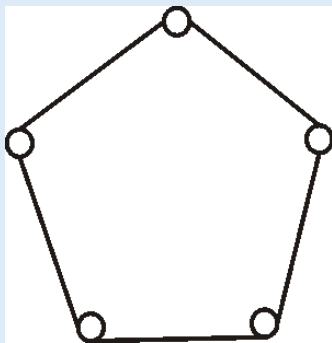
 K_3  K_5 3-vertex directed
complete graph2-regular
undirected graph3-regular graph
Petersen graph4-regular
undirected graph

6.1.3 Common Types of Graphs

↳ Cycle graph and Wheel graph

- **Undirected cycle graph** $C_n = \langle V, E \rangle$, where $V = \{v_1, v_2, \dots, v_n\}$, $E = \{(v_1, v_2), (v_2, v_3), \dots, (v_{n-1}, v_n), (v_n, v_1)\}$, $n \geq 3$, E is a set of unordered pairs.
- **Directed cycle graph** $D_n = \langle V, E \rangle$, where $V = \{v_1, v_2, \dots, v_n\}$, $E = \{\langle v_1, v_2 \rangle, \langle v_2, v_3 \rangle, \dots, \langle v_{n-1}, v_n \rangle, \langle v_n, v_1 \rangle\}$, $n \geq 3$, E is a set of ordered pairs.
- **Wheel graph** W_n : add a single vertex to the undirected cycle graph C_{n-1} and connect it to every vertex of the cycle with exactly one edge, $n \geq 4$.

■ Example:



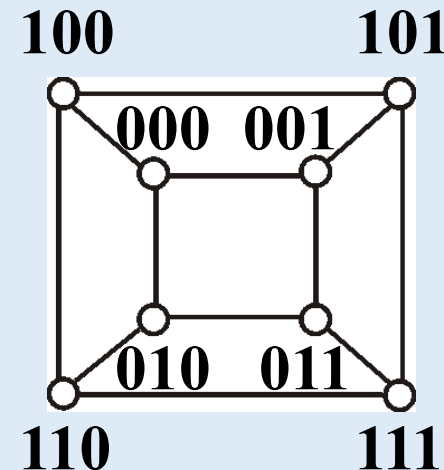
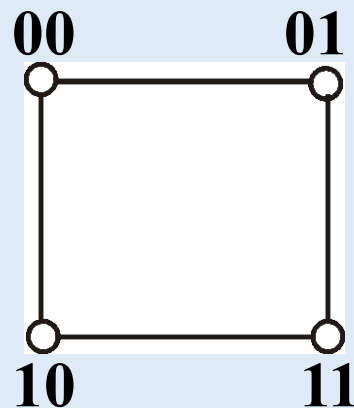
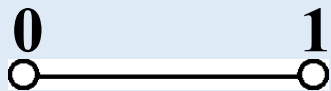
↳ Cycle graph and Wheel graph

- The *n -dimensional hypercube graph* $Q_n = \langle V, E \rangle$ is a simple undirected graph of order 2^n , where

$$V = \{v \mid v = a_1 a_2 \dots a_n, a_i = 0, 1, i = 1, 2, \dots, n\}$$

$$E = \{(u, v) \mid u, v \in V \wedge u \text{ and } v \text{ differ in exactly one bit position}\}.$$

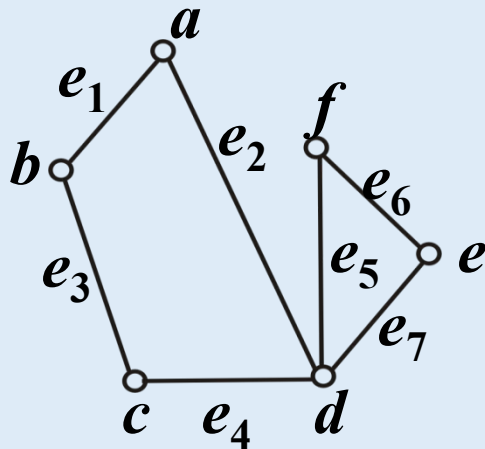
- Example:



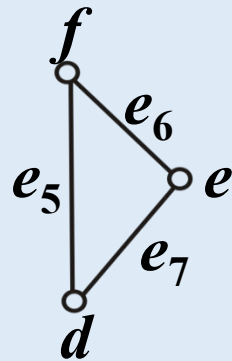
↳ 6.1 Basic Concepts of Graphs

- 6.1.1 Undirected and Directed Graphs
- 6.1.2 Vertex Degree and the Handshaking Lemma
- 6.1.3 Common Types of Graphs
- 6.1.4 Subgraphs and Complements
- 6.1.5 Graph Isomorphism

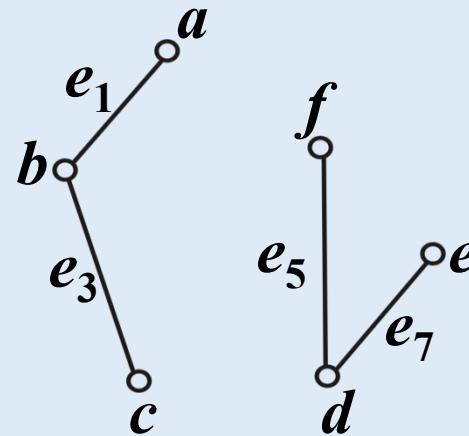
- **Definition 6.4:** Let $G=\langle V,E\rangle$, $G'=\langle V',E'\rangle$ be two graphs (both undirected or both directed).
 - (1) If $V'\subseteq V$ and $E'\subseteq E$, then G' is G 's **subgraph**, G is called G 's **supergraph**, denoted as $G'\subseteq G$.
 - (2) If $G'\subseteq G$ and $V'=V$, then G' is called a **spanning subgraph** of G .
 - (3) If $V'\subset V$ or $E'\subset E$, then G' is called a **proper subgraph** of G .
 - (4) Let $V'\subseteq V$ and $V'\neq\emptyset$, the subgraph of G with vertex set V' and edge set consisting of all edges in G whose endpoints are both in V' is called the **vertex-induced subgraph** on V' , denoted by $G[V']$.
 - (5) Let $E'\subseteq E$ and $E'\neq\emptyset$, the subgraph of G with edge set E' and vertex set consisting of all endpoints of the edges in E' is called the **edge-induced subgraph** on E' , denoted by $G[E']$.



(1)



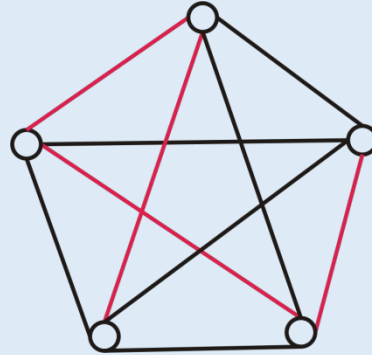
(2)



(3)

- (1),(2),(3) is (1) subgraphs , (2),(3) are proper subgraphs, (1) is the supergraph.
- (1),(3) are spanning subgraphs of (1).
- (2) is the vertex-induced subgraph on $\{d, e, f\}$, and also the edge-induced subgraph on $\{e_5, e_6, e_7\}$.
- (3) is the edge-induced subgraph on $\{e_1, e_3, e_5, e_7\}$.

- **Definition 6.5:** Let $G = \langle V, E \rangle$ be a simple undirected graph of order n , let $\bar{E} = V \times V - E$, then $\bar{G} = \langle V, \bar{E} \rangle$ is called the **complement** of graph G .



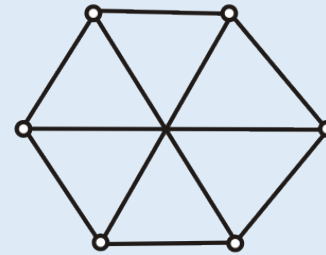
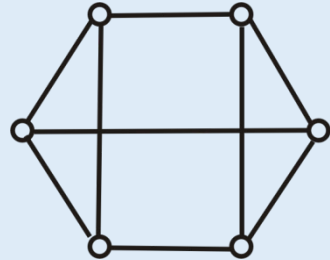
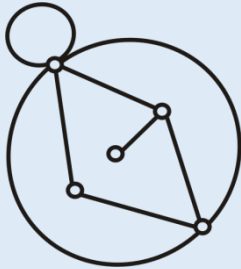
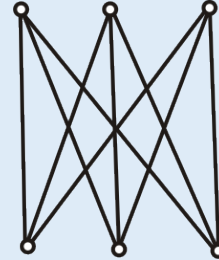
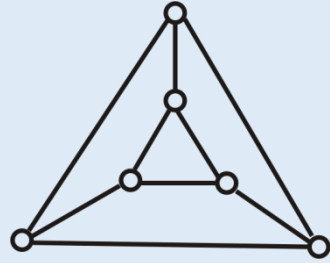
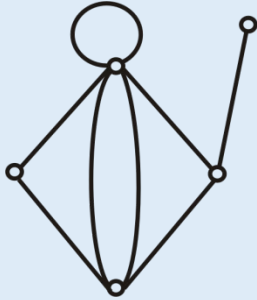
↳ 6.1 Basic Concepts of Graphs

- 6.1.1 Undirected and Directed Graphs
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- 6.1.4 Subgraphs and Complements
- **6.1.5 Graph Isomorphism**

- **Definition 6.6:** Let $G_1 = \langle V_1, E_1 \rangle$, $G_2 = \langle V_2, E_2 \rangle$ be two undirected graphs (or two **directed graphs**), if there exists a bijective function $f: V_1 \rightarrow V_2$, such that for any $v_i, v_j \in V_1$:
- $(v_i, v_j) \in E_1$ if and only if $\langle f(v_i), f(v_j) \rangle \in E_2$
 - and the multiplicity of (v_i, v_j) is equal to $(f(v_i), f(v_j))$
 - then G_1 and G_2 are said to be **isomorphic**, denoted by $G_1 \cong G_2$.

6.1.5 Graph Isomorphism

↳ Isomorphic graphs (e.g.)

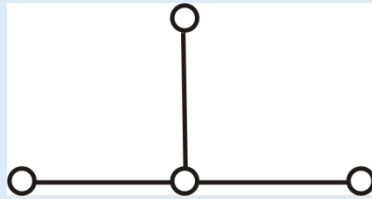


↳ Isomorphic graphs (e.g.)

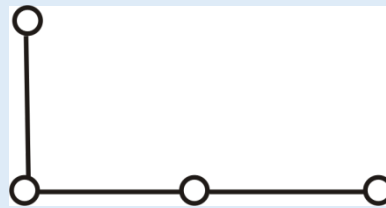
■ **Example** : Draw all non-isomorphic simple undirected graphs with 4 vertices and 3 edges.

■ **Solution**:

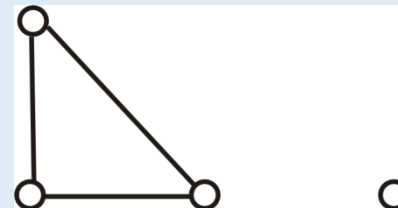
- The total degree is 6 (since the sum of degrees equals twice the number of edges). This must be distributed among 4 vertices.
- The maximum degree is 3, and the number of vertices with odd degree must be even.
- There are three possible degree sequences: **1, 1, 1, 3**; **1, 1, 2, 2**; **0, 2, 2, 2**



1, 1, 1, 3

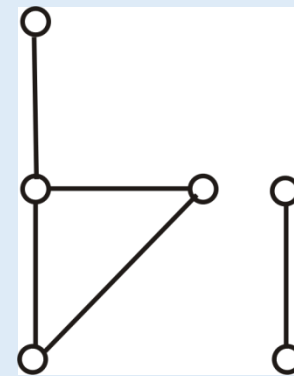
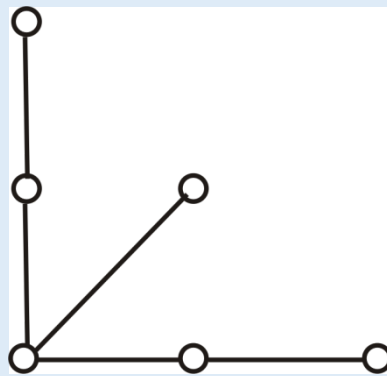
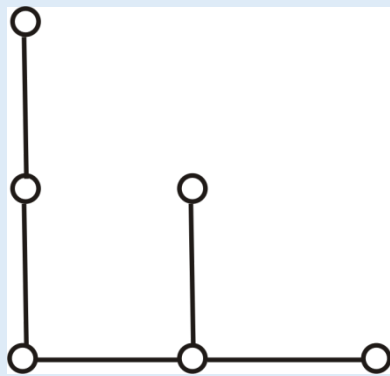


1, 1, 2, 2



0, 2, 2, 2

- **Example:** Draw three non-isomorphic simple undirected graphs with the degree sequence 1,1,1,2,2,3.



6.1 Basic Concepts of Graphs • Brief summary

Objective :

Key Concepts :